# Convex Decompositions of Fuzzy Partitions* 

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#### Abstract

In this paper we investigate some algebraic and geometric properties of fuzzy partition spaces (convex hulls of hard or conventional partition spaces). In particular, we obtain their dimensions, and describe a number of algorithms for effecting convex decompositions. Two of these are easily programmable, and each affords a different insight about data structures suggested by the fuzzy partition decomposed. We also show how the sequence of partitions in any convex decomposition leads to a matrix for which the norm of the corresponding coefficient vector equals a scalar measure of partition fuzziness used with certain fuzzy clustering algorithms.


## I. Introduction and Conclusions

The clustering algorithms of Ruspini [1], Woodbury [2], Bezdek [3], and Dunn [4] all yield fuzzy partitions as clustering solutions for partitioning finite data sets. It was shown in [3] that fuzzy partition spaces are minimal convex supersets, that is, convex hulls, of hard (or conventional) partition spaces. Our goal in this paper is to explore some of the algebraic and geometric consequences of this convexity property.

Partition spaces are defined in Section II; some previous results and several new observations are given. Section III contains our proof that the space of fuzzy $c$-partitions on $n$ data points has dimension $n(c-1)$. In IV an algorithm using a minimax strategy is defined and illustrated numerically. We show that minimax decompositions have lexicographically optimal coefficients. A second decomposition is given in $V$ that interprets the construction of fuzzy partitions from a

[^0]sequence of hard ones as a sequence of transitions or reclassifications from a state of "maximum membership." In VI the partition coefficient used in [5] for evaluation of cluster validity is related to convex decompositions: we prove that this measure of fuzziness is proportional to the number of reclassifications made in the terms of every decomposition of the fuzzy partition in question. This result lends additional support to the clustering strategy proposed in [3] and [5].

## II. Partition Spaces

Let $X=\left\{\mathbf{x}_{1}, \mathbf{x}_{2}, \ldots, \mathbf{x}_{n}\right\} \subset \mathbb{R}^{s}$ be a given finite data set. We fix the integer $c$, $2 \leqslant c<n$, and denote by $V_{c n}$ the usual vector space of real $(c \ll n)$ matrices. Suppose $P=\left(Y_{1}, \mathrm{I}_{2}, \ldots, Y_{c}\right)$ to be a conventional hard $c$-partition of $X$. Thus for each $i, \mathrm{I} \leqslant i \leqslant c$ we have $Y_{\imath} \subset X$; for each $i \neq j$ the intersection $Y_{\imath} \cap Y_{j}$ is empty; and the union of the $Y_{i}$ 's is all of $X, \bigcup_{i=1}^{c} Y_{i}=X$. We say that $P$ is non-degenerate in case none of the $Y_{i}$ 's is empty, and is degenerate otherwise. Partitions of $X$ can be conveniently characterized by matrices in $V_{e n}$ as follows: let $u_{2 k}$ be the $i k$-th element of $U \in V_{r n}$, and define

$$
\begin{equation*}
P_{c}=\left\{U \in V_{c n}: u_{i k} \in\{0,1\} \forall i, k ; \sum_{i=1}^{c} u_{i k}=1 \forall k ; \sum_{k=1}^{n} u_{i k}>0 \forall i\right\} \tag{1}
\end{equation*}
$$

For $U$ in $P_{c}$ we interpret $u_{i k}$ as the value of a characteristic function $u_{i}: X \rightarrow$ $\{0,1\} ; u_{i k}$ specifies the membership of $\mathbf{x}_{k}$ in a partitioning subset $\sum_{\imath}$ of $X$ :

$$
\begin{align*}
u_{2, k} \doteq u_{i}\left(\mathbf{x}_{k}\right) & =1 ; & & \mathbf{x}_{k} \in Y_{i} \\
& =0 ; & & \text { otherwise } \tag{2}
\end{align*}
$$

The $c$-tuple $\left(u_{1}, u_{2}, \ldots, u_{c}\right)$ is the function-theoretic equivalent of ( $Y_{1}, \Sigma_{2}, \ldots$, $Y_{c}$ ), so each $U \in P_{c}$ is uniquely identifiable with a hard $c$-partition of $X$ via (2). Accordingly, we may call $P_{c}$ hard $c$-partition space associated with $X$, and because degeneracy is manifested by zero rows in $U$, the superset

$$
\begin{equation*}
P_{c 0}=\left\{U \in V_{c n}: u_{i k} \in\{0,1\} \forall i, k ; \sum_{i=1}^{c} u_{2 k}=1 \forall k ; \sum_{k=1}^{n} u_{i k} \geqslant 0 \forall i\right\} \tag{3}
\end{equation*}
$$

is degenerate hard c-partition space for $X$.
$P_{c}$ and $P_{c 0}$ are not "spaces" in any ordinary sense; rather, they are extremely large finite sets in the positive or non-negative orthant of $V_{e n}$. Since $\left|P_{c 0}\right|<\infty$, exhaustive search for "optimal" partitions of $X$ is theoretically possible, but
infeasible in practice because of their cardinalities: e.g., $\left|P_{c}\right| \approx 10^{18}$ if $c=10$, $n=25$. Thus finiteness is an impediment to tractable algorithms as well as analytic techniques when partitioning of $X$ is desired.

An even stronger objection to $P_{c}$ lies with the physical interpretation of data substructure it requires. Every $\mathbf{x}_{k}$ in $X$ belongs entirely to one and only one of the hard partitioning subsets $u_{i}$ (since $u_{i}$ and $Y_{i}$ are equivalent, we may call $u_{i}$ a set). All members of $u_{i}$ are fully related to each other, and at the same time totally unrelated to all other members of the data because the boundaries of the partitioning subsets are hard. This is a particulary harsh model for many physical processes, since data representative of most situations originates from mixed populations. It seems more realistic to allow individuals to share memberships in several partitioning subclasses (for example, this is the situation we anticipate for data representing hybrids in mixtures of biological species at the same strata).

A natural way to ameliorate these objections was suggested by Zadeh in [6], who proposed that non-statistical uncertainties of the type described above might be more accurately represented by allowing memberships in fuzzy sets, characterized by membership functions valued in [0, 1]. Motivated by these considerations, and using (1) as our guide, we define

$$
\begin{equation*}
P_{f c}=\left\{U \in V_{c n}: u_{i k} \in[0,1] \forall i, k ; \sum_{i=1}^{c} u_{i k}=1 \forall k ; \sum_{k=1}^{n} u_{i k}>0 \forall i\right\} \tag{4}
\end{equation*}
$$

as fuzzy, non-degenerate c-partition space for $X$. Here $u_{i k}$ is again the grade of membership of $\mathbf{x}_{k}$ in the fuzzy subset $u_{i}: X \rightarrow[0,1]$. The condition $\sum_{i=1}^{c} u_{i k}=1$ for each $k$ insures that each $\mathbf{x}_{k}$ has unit membership (in $X$ ); these memberships may be distributed among the $c$ fuzzy subsets $\left\{u_{i}\right\}$ arbitrarily as long as their sum is unity. Corresponding to $P_{c 0}$ in the hard case is degenerate fuzzy c-partition space $P_{\text {fe0 }}$ obtained from (4) by relaxing the last condition exactly as was done in (3); $P_{f c 0}$ is not used in the sequel.

A substantial amount of information concerning the imbeddings $P_{c} \subset P_{c 0} \subset$ $P_{f c} \subset P_{f c 0}$ is available elsewhere [3]. The main fact established there we intend to exploit below is that $P_{f 0}$ is the convex hull of $P_{c 0}, P_{f c}=\operatorname{conv}\left(P_{c 0}\right)$. We observe that the convex hull of $P_{c}$ is a proper subset of $\operatorname{conv}\left(P_{c 0}\right)$, by noting that, for example, with any $\lambda$ in $[0,1]$ the matrix

$$
U=\left[\begin{array}{ccc}
\lambda & \lambda & \lambda  \tag{5}\\
1-\lambda & 1-\lambda & 1-\lambda
\end{array}\right]
$$

lies in $\operatorname{conv}\left(P_{20}\right)$, but is not in $\operatorname{conv}\left(P_{2}\right)$. In other words, this $U$ has no convex decomposition with all non-degenerate terms. The additional property $U$ in $P_{f c}$ needs to distinguish it as a member of $\operatorname{conv}\left(P_{c}\right)$ is not yet known: our interest lies with $P_{f c}$ due to the physical considerations outlined above.

## III. The Dimension of Fuzzy Partition Space

In this section we prove that the dimension of $P_{f c}$ is $n(c-1), \operatorname{dim}\left(P_{f c}\right)=$ $n(c-1) . P_{f c}$ is convex, so takes its dimension from the vector subspace which translates the affine hull of $P_{f c}$ to the origin of $V_{c n} ; \operatorname{dim}\left(P_{f c}\right)=\operatorname{dim}(M)$, where $\operatorname{aff}\left(P_{f c}\right)=U^{*}+M, U^{*} \in V_{c n}$, and $M$ a vector subspace in $V_{c n}$. Before stating and proving this theorem, we sketch the idea of the proof. Given $U \in P_{f r}$, one can choose any positive path $\mathbf{v}_{3}=\left(*_{1}, *_{2}, \ldots, *_{n}\right)$ with $*_{2}>0$ for every $i$ through the columns of $U$. Such a path is illustrated pictorially in (6):


Call $c_{j}$ the smallest element of $\mathbf{v}_{j}$; define a matrix $U_{j} \in P_{c 0}$ to have 1's at every address of the entries of $\mathbf{v}_{j}$, and 0 's elsewhere. Define a residual matrix $R_{j}=$ $U-c_{j} U_{j}$. Apply this iteratively to $R_{j}$, beginning at $R_{0} \doteq U$; we shall first show that $R_{j} \rightarrow \theta$ in $n(c-1)+1$ or less steps; and to complete the proof, that there is a matrix in $P_{f c}$ which cannot be decomposed with less than $n(c-1)+1$ terms. This will prove

Theorem 1. For $P_{f c}$ as defined in (5), we have

$$
\begin{equation*}
\operatorname{dim}\left(P_{f c}\right)=n(c-1) \tag{7}
\end{equation*}
$$

Proof. Let $U \in P_{f c}$. For $1 \leqslant k \leqslant n$ choose $\boldsymbol{i}_{k}$ so that the entries $u_{i_{k}, k}>0$, and set $L^{\zeta}=R_{0}, I_{1}=\left\{i_{k}: 1 \leqslant k \leqslant n\right\}$. This is always possible because columns of $U$ must sum to 1 , so every column has at least one non-zero entry. Define the $n$-vector $\mathbf{v}_{1}=\left(u_{i_{1}, 1}, u_{i_{2}, 2}, \ldots, u_{i_{n}, n}\right)$, and set $c_{1}=\min _{l_{1}}\left\{u_{i_{k}, k}\right\}$. If $c_{1}=1, U$ is already hard and we are done. Otherwise, define the matrix $U_{1} \in P_{c 0}$ via

$$
\begin{array}{rlrl}
\left(u_{1}\right)_{i_{k}, k} & =1 ; & & \\
& i_{k} \in I_{1}  \tag{8}\\
& =0 ; & & \text { otherwise }
\end{array}
$$

so that $U_{1}$ has l's wherever $\mathbf{v}_{1}$ passes through $U$, and 0 's elsewhere. Next, define the residual matrix

$$
\begin{equation*}
R_{1}=U-c_{1} U_{1}=R_{0}-c_{1} U_{1} \tag{9}
\end{equation*}
$$

Note that every column of $R_{1}$ sums to $1-c_{1}>0$; moreover, $R_{1}$ contains at least one zero, viz., at the address where $c_{1}$ occured in $U$ (it may have more than one zero if $\min _{I_{1}}\left\{u_{i_{k_{k}}, k}\right\}$ is not unique).

Proceeding iteratively, we apply the steps above to $R_{1}$ : find a non-zero path $\mathbf{v}_{2}$ through its columns, factor out the smallest element - $c_{2}$ - in the path; define $U_{2} \in P_{c 0}$ using these path addresses, and set $R_{2}=R_{1}-c_{2} U_{2}$. Either $R_{2}$ is the zero matrix and we are done, or (i) column sums of $R_{2}$ are equal to $1-\left(c_{1}+c_{2}\right)$ $>0$, and (ii) $R_{2}$ has at least two zeores.

Continuing inductively, we have at the $j$-th step the residual $R_{j}=U$ $\sum_{k=1}^{j} c_{k} U_{k}$, and if $R_{j}$ is not the zero matrix, then (i) column sums of $R_{j}$ are positive, $\sum_{i=1}^{e}\left(r_{j}\right)_{i k}=1-\sum_{k=1}^{j} c_{k}$, and (ii) $R_{j}$ has at least $j$ zeroes. Now let $m=n(c-1)$. Either this process terminates at some $j \leqslant m$, or after $m$ iterations the residual $R_{m}$ has at least $m$ zeroes. If $R_{m}$ is not the zero matrix, its positive column sums are $1-\sum_{k=1}^{m} c_{k}>0$, so every column has a non-zero entry. But $m=n(c-1)$, so each column of $R_{m}$ has exactly one entry which must equal $1-\sum_{k=1}^{m} c_{k}$. Define $c_{m+1}=1-\sum_{k=1}^{m} c_{k}$, and $U_{m+1} \in P_{c 0}$ using the addresses of the entries in $R_{m}$ in the usual way. Then $R_{m+1}=U-\sum_{k=1}^{m+1} c_{k} U_{k}=$ the zero matrix. Since $\sum_{k=1}^{m+1} c_{k}=1$ with all of the $c_{k}$ 's in $(0,1)$ and all of the $U_{k}$ 's $\in P_{c 0}$, the convex decomposition of $U$ is complete.

We have shown that every $U \in P_{f c}$ admits at least one convex decomposition with $n(c-1)+1$ terms. Thus by Caratheodory's theorem for convex sets (cf. Roberts and Varberg [7]), we conclude that $\operatorname{dim}\left(P_{f c}\right) \leqslant n(c-1)$.

To complete the proof we must show that equality prevails. Towards this end, let $k \in \mathbb{R}$ and define $U=\left[u_{i j}\right]$ by

$$
\begin{align*}
u_{i j} & =\frac{k}{(2 m)^{n(i-1)+(j-1)}} ; & & 1 \leqslant i \leqslant c-1 \quad \text { and } \quad 1 \leqslant j \leqslant n  \tag{10}\\
& =1-\sum_{s=1}^{c-1} u_{s j} ; & & i=c \quad \text { and } \quad 1 \leqslant j \leqslant n
\end{align*}
$$

where again $m=n(c-1)$. For $k \in(0,1)$ this matrix belongs to $P_{f c}$. Suppose this $U$ can be decomposed in less than or equal to $m$ terms, say for convenience that $U=\sum_{j=1}^{n} \sum_{i=1}^{c-1} c_{i j} U_{i j}$, where $U_{i j} \in P_{c 0}, c_{i j} \in[0,1] \forall i, j$, and $\sum_{j=1}^{n} \sum_{i=1}^{c-1} c_{i j}$ $=1$. To write $u_{11}=k$ as a partial sum of the convex coefficients, it is necessary that at least one of the $c_{i j}$ 's is contained in the closed interval [ $\left.(1 / m) k, k\right]$. We can assume without loss that $c_{11}$ is this coefficient. Similarly, there is for $u_{12}$ at least one of the $c_{i 1}$ 's, say $c_{12}$, so that $c_{12} \in[(1 / m)(k / 2 m),(k / 2 m)]$, and $c_{11} \neq c_{12}$, because

$$
\left[\left(\frac{1}{m}\right)(k), k\right] \cap\left[\left(\frac{1}{m}\right)\left(\frac{k}{2 m}\right),\left(\frac{k}{2 m}\right)\right]-\varnothing \quad \text { for } k>0
$$

Continuing inductively, it is clear that each $\mathrm{u}_{1 j}$ in (10) for $1 \leqslant i \leqslant c-1$ and $1 \leqslant j \leqslant n$ requires a distinct coefficient $c_{2 j}$ which satisfies

$$
\begin{align*}
& c_{2 j} \in\left[\left(\frac{1}{m}\right)\left(\frac{k}{(2 m)^{n(2-1)+(\jmath-1)}}\right),\left(\frac{k}{(2 m)^{n(i-1)+(\jmath-1)}}\right)\right] ;  \tag{1la}\\
& c_{\imath \jmath} \neq c_{s t} \quad \text { unless } \quad i=s \quad \text { and } \quad j=t . \tag{1lb}
\end{align*}
$$

In view of (11a), we can bound above the sum of these $m c_{i j}$ 's by summing the upper endpoints of the intervals at (11a):

$$
\begin{aligned}
\sum_{j=1}^{n} \sum_{i=1}^{c-1} c_{\imath \jmath} & \leqslant \sum_{j=1}^{n} \sum_{l=1}^{c-1}\left\{\frac{k}{\left.(2 m)^{n(2-1)+(\jmath-1)}\right\}}\right. \\
& =k\left\{1+\frac{1}{(2 m)}+\frac{1}{(2 m)^{2}}+\cdots+\frac{1}{(2 m)^{m-1}}\right\} \\
& =k\left\{\frac{1-\left(\frac{1}{2 m}\right)^{m}}{1-\left(\frac{1}{2 m}\right)}\right\} .
\end{aligned}
$$

If we choose $k$ so that

$$
\begin{equation*}
k<\left\{\frac{1-\left(\frac{1}{2 m}\right)}{1-\left(\frac{1}{2 m}\right)^{m}}\right\} ; \quad m=n(c-1) \tag{12}
\end{equation*}
$$

then $\sum_{j=1}^{n} \sum_{i=1}^{c-1} c_{\imath j}<1$, and we nced at least onc more coefficient to effect a convex decomposition of $U$. That is, we must use $m+1$ terms for the matrix at (10). Given $c$ and $n$, it is clear from (12) that $k$ can always be chosen so that $U$ in (10) belongs to $P_{f c}$, so the proof is complete.
Q.E.D.

## IV. Minimax Decomposition

Because the residual at each step in the algorithm used for the proof of Theorem 1 has at least one entry forced to zero, we call it a forcing (or F) algorithm. Not all convex decompositions are achieved with forcing algorithms, and we shall see in the next section that every $U \in P_{f,}$ can be decomposed without using an $F$ algorithm. Many different $F$-decompositions of a given $U$ will generally be possible, since the only constraint on the path $\mathbf{v}_{j}$ in the above proof was positivity. Despite this problem of non-uniqueness, there is among the $F$-decompositions of $U$ one which deserves special attention. If at each step the
path $\mathbf{v}_{j}$ chosen through $R_{j}$ is one of maximum residual memberships, the $c_{j}$ factored from $\mathbf{v}_{j}$ will be the minimum maximum residual membership, giving rise to the Minimax (or MM) Decomposition.

Let $U \in P_{f c}$. Following the algorithm given in the proof of Theorem 1, we modify the choice of path $\mathbf{v}_{j}$ as follows: at step $j+1$, choose $\mathbf{v}_{j+1}$ through $R_{j}$ by letting $\left(i_{k}, k\right)$ be the index at which $\max _{1 \leqslant i \leqslant 0}\left\{\left(r_{j}\right)_{i k}\right\}$ is achieved for $1 \leqslant k \leqslant n$. If more than one index occurs at any $k$, choose any one of these. Set $\mathbf{v}_{j+1}=$ $\left(*_{i_{1}, 1}, *_{i_{2}, 2}, \ldots, *_{i_{n}, n}\right)$, and $c_{j+1}=\min _{1 \leqslant k \leqslant n}\left\{\left(r_{j}\right)_{i_{k}, k}\right\}$. Then define $U_{j+1} \in P_{c 0}$ to have l's at the addresses from which the entries of $\mathbf{v}_{j+1}$ were chosen, and 0 's elsewhere. Finally, put $R_{j+1}=R_{j}-c_{j+1} U_{j+1}$, and iterate until $R_{k} \rightarrow \theta$, the zero matrix.

To illustrate $F$-decompositions, we apply first an arbitrary $F$-algorithm, and then the $M M$-algorithm to the matrix

$$
U=\left[\begin{array}{ccccc}
.90 & .80 & .30 & .40 & .05  \tag{13}\\
.10 & .20 & .70 & .60 & .95
\end{array}\right]
$$

## (i) Arbitrary $F$-decomposition

Since $u_{i k}$ is positive for every $i$ and $k$, we can extract at once any $c_{1}=u_{i k}$. Arbitrarily choosing path $\mathbf{v}_{1}=\left(*_{21}, *_{22}, *_{23}, *_{24}, *_{15}\right)$ and taking for $c_{1}$ entry $u_{15}=.05$ leads to

$$
c_{1} U_{1}=.05\left[\begin{array}{lllll}
0 & 0 & 0 & 0 & 1 \\
1 & 1 & 1 & 1 & 0
\end{array}\right]
$$

and

$$
R_{1}=U-c_{1} U_{1}=\left[\begin{array}{ccccc}
.90 & .80 & .30 & .40 & 0 \\
.05 & .15 & .65 & .55 & .95
\end{array}\right]
$$

Choosing $c_{2}=\left(r_{1}\right)_{21}=.05$ with $\mathbf{v}_{2}=\left(*_{21}, *_{22}, *_{23}, *_{24}, *_{25}\right)$ now yields

$$
c_{2} U_{2}=.05\left[\begin{array}{lllll}
0 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 1 & 1
\end{array}\right]
$$

and

$$
R_{2}=\left[\begin{array}{ccccc}
.90 & .80 & .30 & 40 & 0 \\
0 & .10 & .60 & .50 & .90
\end{array}\right]
$$

Continuing in this fashion, one may generate various $F$-decompositions of $U$. For example, the paths and coefficients

$$
\begin{array}{ll}
c_{3}=.10: & \mathbf{v}_{3}=\left(*_{11}, *_{22}, *_{23}, *_{24}, *_{25}\right) \\
c_{4}=.30: & \mathbf{v}_{4}=\left(*_{11}, *_{12}, *_{13}, *_{14}, *_{25}\right) \\
c_{5}=.10: & \mathbf{v}_{5}=\left(*_{11}, *_{12}, *_{23}, *_{14}, *_{25}\right) \\
c_{6}=.40: & \mathbf{v}_{6}=\left(*_{11}, *_{12}, *_{23}, *_{24}, *_{25}\right)
\end{array}
$$

result in the particular decomposition

$$
\begin{align*}
& {\left[\begin{array}{ccccc}
.90 & .80 & .30 & .40 & .05 \\
.10 & .20 & .70 & .60 & .95
\end{array}\right]} \\
& \quad=0.5\left[\begin{array}{lllll}
0 & 0 & 0 & 0 & 1 \\
1 & 1 & 1 & 1 & 0
\end{array}\right] \quad .05\left[\begin{array}{lllll}
0 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 1 & 1
\end{array}\right]+.10\left[\begin{array}{lllll}
1 & 0 & 0 & 0 & 0 \\
0 & 1 & 1 & 1 & 1
\end{array}\right] \\
& \quad-.30\left[\begin{array}{lllll}
1 & 1 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 1
\end{array}\right]+.10\left[\begin{array}{lllll}
1 & 1 & 0 & 1 & 0 \\
0 & 0 & 1 & 0 & 1
\end{array}\right]+.40\left[\begin{array}{lllll}
1 & 1 & 0 & 0 & 0 \\
0 & 0 & 1 & 1 & 1
\end{array}\right] . \tag{14}
\end{align*}
$$

In (14) only ${C_{2}}_{2}$ is degenerate: the proportion of degenerate terms appearing in $F$-decompositions tends to increase with $c$. Moreover, there is little practical value for a decomposition like (14), since it furnishes no explanation or interpretation for structure in the fuzzy partition it decomposes.

## (ii) MM-decomposition

Beginning again with $U$ in (13), we find the initial minimax path $\mathbf{v}_{1}=\left({ }_{11},{ }_{12}, *_{23}, *_{24}, *_{25}\right)$ through $U$ : the address yielding the $M M$ membership is ( 2,4 ), so $c_{1}=u_{24}=.60$, thus

$$
c_{1} U_{1}==.60\left[\begin{array}{lllll}
1 & 1 & 0 & 0 & 0 \\
0 & 0 & 1 & 1 & 1
\end{array}\right]
$$

and

$$
R_{1}=R_{0}-c_{1} U_{1}=U-c_{1} U_{1}=\left[\begin{array}{ccccc}
.30 & .20 & .30 & .40 & .05 \\
.10 & .20 & .10 & 0 & .35
\end{array}\right]
$$

Inspection of $R_{1}$ reveals that the residual $M M$ membership is .20 , that is, $c_{2}$ will be .20. There are, however, two $M M$ paths which are available for this choice;

$$
\mathbf{v}_{2}=\left(*_{11}, *_{12}, *_{13}, *_{14}, *_{25}\right)
$$

and

$$
\mathbf{v}_{2}^{\prime}=\left(*_{11}, *_{22}, *_{13}, *_{14}, *_{25}\right)
$$

The choice of path is arbitrary, leading to either one of

$$
c_{2} C_{2}=.20\left[\begin{array}{lllll}
1 & 1 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 1
\end{array}\right] \quad \text { or } \quad c_{2} U_{2}^{\prime}=.20\left[\begin{array}{lllll}
1 & 0 & 1 & 1 & 0 \\
0 & 1 & 0 & 0 & 1
\end{array}\right] .
$$

These choices would yield for residuals the matrices
$R_{2}=\left[\begin{array}{ccccc}.10 & 0 & .10 & .20 & .05 \\ .10 & .20 & .10 & 0 & .15\end{array}\right] \quad$ or $\quad R_{2}^{\prime}=\left[\begin{array}{ccccc}.10 & .20 & .10 & .20 & .05 \\ .10 & 0 & .10 & 0 & 15\end{array}\right]$.

In either case, note that $c_{3}$ will be .10 and at step 3 there will be four $M M$ paths available, two for each choice for $R_{2}$. All of these choices will leave residuals with minimax membership in $R_{3}$ equal to $c_{4}=.05$. Finally, every branch at the fifth step would require the choice $c_{5}=.05$. We exhibit one $M M$ decomposition of $U$ :

$$
\begin{align*}
& {\left[\begin{array}{ccccc}
.90 & .80 & .30 & .40 & .05 \\
.10 & .20 & .70 & .60 & .95
\end{array}\right] } \\
& \quad=.60\left[\begin{array}{lllll}
1 & 1 & 0 & 0 & 0 \\
0 & 0 & 1 & 1 & 1
\end{array}\right]+.20\left[\begin{array}{lllll}
1 & 1 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 1
\end{array}\right]+.10\left[\begin{array}{lllll}
0 & 0 & 0 & 1 & 0 \\
1 & 1 & 1 & 0 & 1
\end{array}\right] \\
& \quad .05\left[\begin{array}{lllll}
1 & 0 & 1 & 1 & 1 \\
0 & 1 & 0 & 0 & 0
\end{array}\right]+.05\left[\begin{array}{lllll}
1 & 0 & 1 & 1 & 0 \\
0 & 1 & 0 & 0 & 1
\end{array}\right] . \tag{15}
\end{align*}
$$

Although the $M M$-decomposition (15) is non-unique, there are several important aspects of it that are invariant. In what follows it is convenient to introduce the following terminology:

Definition 1. Let $\sum_{k=1}^{P} c_{k} U_{k}$ be any convex decomposition of $U$ in $P_{f c}$. We call $p$ the length of the decomposition, and the hard partition $U_{3} \in P_{c 0}$ with $c_{j}=\max _{1 \leqslant k \leqslant p}\left\{c_{k}\right\}$ is called its dominant term. Since the sequence $\left\{c_{k}\right\}$ of coefficients generated by $M M$-algorithms is monotone decreasing by their definition, i.e., $1 \geqslant c_{1} \geqslant c_{2} \cdots \geqslant c_{p}>0, U_{1}$ is always the dominant term in $M M$ decomposition. In particular, this $U_{1}$ is the hard partition of $X$ "closest" to $U$ in the sense of maximum membership (cf. [5]), and deserves to be isolated by a special notation, say $U_{m m}$, because it is one whose structure can be readily interpreted as the best (in the sense of maximum membership) hard approximation to $U \in P_{f c}$.

Theorem 1 shows that the $M M$-algorithm terminates in at most $n(c-1)+1$ steps. In general, however, it may stop sooner, as was the case in our example above. Roughly speaking, this happens because we extract at each step the largest possible coefficient, thus yielding decompositions of minimal length. While $M M$ decompositions are often non-unique, we may prove that both their lengths and coefficient sequences are:

Theurem 2. If $U \in P_{f c}$ has two distinct MM-decompositions, say $U=$ $\sum_{k=1}^{p} c_{k} U_{k}=\sum_{k=1}^{q} c_{k}^{\prime} U_{k}^{\prime}$, then

$$
\begin{equation*}
p=\boldsymbol{q} \tag{16a}
\end{equation*}
$$

and

$$
\begin{equation*}
c_{k}=c_{k}^{\prime} \quad \forall 1 \leqslant k \leqslant p \tag{16b}
\end{equation*}
$$

Proof. Let $M$ and $N$ belong to $V_{c n}$, and set $I_{c}=\{1,2, \ldots, c\}$. We shall say that $M \cdot c p \cdot N$ iff for $1 \leqslant j \leqslant n$ there is a permutation $\pi_{j}: I_{c} \rightarrow I_{c}$ so that $m_{1 j}=n_{\pi_{2}(d), j}, 1 \leqslant i \leqslant c$. In other words, $M$ and $N$ can be obtained from each other by independently permuting the entries in corresponding columns.

Now let $U \in P_{f c}$ and suppose at the $k$-th step in $M M$-decomposition of $U^{-}$ two distinct $M M$ paths $\mathbf{v}_{k}$ and $\mathbf{v}_{k}^{\prime}$ through $R_{k-1}$ are available, yielding hard partitions $U_{k} \not \boldsymbol{z}^{\prime} U_{k}^{\prime}$. The values of $\mathbf{v}_{k}$ and $\mathbf{v}_{k}^{\prime}$ are equal: it is their address origins that differ at least once, leaving for residuals the matrices $R_{k} \neq R_{h}^{\prime}$. Since the columns of $R_{k}$ and $R_{k}^{\prime}$ are identical up to a permutation of each pair of columns at addresses where $\mathbf{v}_{k}$ and $\mathbf{v}_{k}^{\prime}$ differ, $R_{k} \cdot c p \cdot R_{k}^{\prime}$.

Applying the $M M$-algorithm to $R_{k}$ and $R_{k}^{\prime}$, suppose $c_{k+1}$ and $c_{h+1}^{\prime}$ to be the coefficients extracted, respectively. $R_{k} \cdot c p \cdot R_{k}^{\prime}$ implies that all column maximums in $R_{k}$ and $R_{k}^{\prime}$ are equal, so $c_{k+1}=c_{k+1}^{\prime}$. Repetition of the argument above now yields for the new residuals $R_{k+1} \cdot c p \cdot R_{k+1}^{\prime}$. In view of this, identical coefficients must be extracted at every step. From which there follows (16a) and (16b).
Q.E.D.

Now suppose that $\sum_{k=1}^{q} d_{k} W_{k}$ is any convex decomposition of $U \in P_{f c}$, and that $\left\{c_{1}, c_{2}, \ldots, c_{p}\right\}$ are the coefficients for a $M M$-decomposition. Assuming without loss that the $\left\{d_{k}\right\}$ are monotone decreasing, $1 \geqslant d_{1} \equiv d_{2} \cdots \geqslant d_{q}>0$, we note that $c_{1} \geqslant d_{1}$, because $c_{1}$ is the largest possible coefficient that can be factored from $U$. If $c_{1}=d_{1}$ we may take $W_{1}=U_{1}$ in the $M M$-decomposition, i.e., $W_{1}-U_{1}=U_{m m}$. Continuing in this fashion, we can construct from the $W_{k}$ 's a $M M$ sequence as long as the $d_{k}$ 's continue to equal the $c_{k}$ 's. Either this process terminates with $p=q$ and $d_{k}=c_{k} \forall k$; or at some $j, d_{j}<c_{1}$. This proves

Theorem 3. Let $U \in P_{f c}$. The coefficient vector $\mathrm{c}_{n, m}=\left(c_{1}, c_{2}, \ldots, c_{p}\right)$ for all MM-decompositions of $U$ is lexicographically larger than the coefficient vector $\mathbf{d}=\left(d_{1}, d_{2}, \ldots, d_{q}\right)$ for any other convex decomposition of $L$.

Since the sum of convex coefficients in every convex decomposition of $C^{r}$ is one, it is plausible to presume that the size of each coefficient is a measure of the relative efficacy of its associated hard partition of $X$ as a factor in the construction of fuzzy partition $U$. Theorem 3 shows that $U_{m i m}$, the first term of every $M M$ decomposition of $U$, is always the "strongest" possible dominant term, and the remaining terms in $M M$-decompositions form the optimal sequence of factor partitions in the sense of coefficient magnitudes. Empirical evidence as well as one's intuition suggest that a related property of $M M$-decomposition is also true:

Conjecture. Let $U \in P_{f c}$. If $\sum_{k=1}^{p} c_{k} U_{k}$ is any $M M$-decomposition of $U$, and $\sum_{k=1}^{q} d_{k} W_{k}$ is any other decomposition of $U$, then $p \leqslant q$.

Indeed, $p=2$ implies $q \geqslant 2$, for otherwise $q=1$ implies $U \in P_{c 0}$ so $p=1$. At $p=3$ the same argument shows that $q \neq 1$, so the conjecture is true unless
$q=2$. If $q=2$, then $c_{1}+c_{2}+c_{3}=d_{1}+d_{2}=1$. Theorem 3 insures us that $c_{1} \geqslant d_{1}$ and $c_{1} \geqslant d_{2}$. If $c_{1}$ were equal to either $d_{1}$ or $d_{2}$, we could choose for the corresponding $W_{j}$ the matrix $U_{1}$, leaving equal residuals $U-c_{1} U_{1}=U-d_{j} W_{j}$. Then the remaining $d_{k}$ must equal $c_{2}$, for otherwise maximality of $c_{2}$ contradicts $c_{2} \geqslant d_{k}$, and $d_{k} \leqslant c_{2}$ implies that $1=d_{j}+d_{k}, c_{1}+c_{2}=1$. In either case, this leaves $c_{3}=0$ so $p \neq 3$. On the other hand, if $c_{1}>d_{1}$ and $c_{1}>d_{2}$, then because $c_{1}=u_{i j}$ for some $i$ and $j$, it is necessary that $c_{1}=d_{1}+d_{2}=1$, so $c_{2}=c_{3}=0$, and again $p \neq 3$, thus, $p=3$ implies $q \geqslant 3$. These remarks establish the conjecture for $p \leqslant 3$, and strongly support our guess that $M M$-decomposition is always minimal in length.

## V. Reclassification Decomposition

$F$-decompositions force at least one residual membership to zero at each step in their application: we interpret this as forcing the individual concerned out of its membership in the class it leaves. In $M M$-decomposition all $n$ subjects begin in their maximum membership classes (in $U_{1}=U_{n: m}$ ), and when $c_{2}$ is extracted as many individuals as possible are allowed to shift classes in the hard partition $U_{2}$. An algorithm for decomposition which is based on membership thresholding and proceeds quite differently from the $M M$-algorithm is the

## Reclassification (or $R$ ) Decomposition

Let $U \in P_{f c}$. From $U$ we first construct a ranking matrix $\rho(U) \in V_{c n}$ as follows: $\rho_{i j}$ is the integer in $I_{c}=\{1,2, \ldots, c\}$ which gives the rank of $u_{i j}$ in column $j, 1 \leqslant j \leqslant n$, after ordering the elements of each column of $U$ in descending order.

A second matrix $\sigma(U) \in V_{c n}$ is then derived from $\rho(U)$ by taking cumulative sums of entries in $U$ : wherever $\rho_{i j}=1$, define $\sigma_{i j}=u_{i j}$; where $\rho_{i j}=2$, define $\sigma_{2 j}=u_{i_{1}, j}+u_{i j}$, where $\rho_{i_{1}, j}=1$; and so on. In other words, the entries in columms of $\sigma(U)$ are cumulative sums of the memberships of each $\mathbf{x}_{k}$ in the $c$ fuzzy sets exhibited in $U$ : the values in column $j$ ascend from the largest memberships, and so on up to 1 , the addresses of this ordered sequence being found at $1,2, \ldots, c$ in the corresponding column of $\rho(U)$. For example, with

$$
U=\left[\begin{array}{llll}
.70 & .40 & .03 & .15  \tag{18}\\
.20 & .50 & .07 & .80 \\
.10 & .10 & .90 & .05
\end{array}\right]
$$

we have

$$
\rho(U)=\left[\begin{array}{llll}
1 & 2 & 3 & 2 \\
2 & 1 & 2 & 1 \\
3 & 3 & 1 & 3
\end{array}\right] ; \quad \sigma(U)=\left[\begin{array}{cccc}
.70 & .90 & 1 & .95 \\
.90 & .50 & .97 & .80 \\
1 & 1 & .90 & 1
\end{array}\right]
$$

Next, we array the $n(c-1)$ smallest positive entries of $\sigma(U)$ in increasing order, label them $a_{1}$ to $a_{m}, m=n(c-1)$, define $a_{0}=0 ; a_{m+1}=1$, and calculate their successive differences:

$$
\begin{gather*}
0=a_{0}<a_{1} \leqslant a_{2}, \quad \cdots \leqslant a_{m} \leqslant a_{m+1}=1  \tag{19}\\
c_{j+1}=a_{j+1}-a_{j} ; \quad 0 \leqslant j \leqslant m . \tag{20}
\end{gather*}
$$

The $c_{j}$ 's all lie in [0,1], and by their definition sum to $1, \sum_{l=1}^{m+1} c_{j}=a_{m+1}-a_{0}$ $=1$ : these are the convex coefficients of the $R$-decomposition of $U$. Since $a_{1}$ is the smallest positive entry in $\sigma(U), c_{1}=a_{1}-a_{0}$ is the minimum maximum membership in $U$. That is, the first coefficient extracted by both the $R$ and MM algorithms is always the same one - the largest possible one. For the matrix at (18) we have from ( $U$ )

$$
\begin{align*}
& a_{0}=.00  \tag{21}\\
& a_{1}=.50>c_{1}=.50 \\
& a_{2}=.70>c_{2}=.20 \\
& a_{3}-.80>c_{3}=.10 \\
& a_{4}=.90>c_{4}=.10 \\
& a_{5}=.95>c_{5}=.05 \\
& a_{6}=.97>c_{6}=.02 \\
& a_{7}=1
\end{align*} c_{7}=.03 .
$$

Note that cven though $c_{1}$ is always the largest coefficient extracted, this sequence is not necessarily monotone decreasing; it will be unique and of length $n(c-1)+1$ whenever the ordering in (19) is strict. To complete the decomposition, we define hard partitions $U_{k} \in P_{c 0}$ corresponding to the $c_{k}$ 's as follows: $U_{1}=U_{m m}$ is again the maximum membership matrix, and is the dominant term of both the $R$ and $M M$ decompositions; every $\mathbf{x}_{k} \in X$ belongs initially to its class of maximum membership. Now think of $r_{k}=\sum_{j=1}^{k} c_{j}$ as the threshold at step $k$ in the $R$-algorithm; note that $t_{k}=a_{k}$ for $k>1$. At step $k$ locate the address (or addresses) in $\sigma(U)$ where threshold $a_{k}$ occurs, and define $U_{k}$ equal to $U_{k-1}$ except: at each address where $a_{k}$ occurs in $\sigma(U)$, set $\left(u_{k}\right)_{i j}=0$, and move the I's that occupied these addresses to the place where their next lowest membership occurs, i.e., to the next highest integer address in the corresponding column of ranking matrix $\rho(U)$. This transition rule amounts to reclassifying between $U_{\text {, }}$ and $U_{j}, i<j$, all of those vectors which are thresholded from their upper to lower "halves" in $\sigma(U)$ as theshold $t$ runs through $\left(t_{1}, t_{1}\right]$. This procedure is admittedly hard to described verbally, but should be clear from the example to follow: from (18) and (21) we would proceed by imagining threshold $t$ to begin at 0 ; for $0 \leqslant t \leqslant .50$ no reclassification is made, and we have

$$
c_{1} U_{1}=c_{1} U_{m m}=.50\left[\begin{array}{llll}
1 & 0 & 0 & 0  \tag{22}\\
0 & 1 & 0 & 1 \\
0 & 0 & 1 & 0
\end{array}\right] .
$$

As soon as $t$ becomes greater than .50 , vector $\mathbf{x}_{2}$ cannot tolerate membership in class 2 (one sees this by finding $t=.50$ at address $(2,2)$ in $\sigma(U)$ ); to find the destination of $\mathbf{x}_{2}$, go to address $(2,2)$ in $\rho(U)$, and find the address of the next highest integer, in this case, address (1,2). Thus the transition from $U_{1}$ to $U_{2}$ at $t=0.5$ is made by sending $x_{2}$ from $(2,2)$ in $U_{1}$ to $(1,2)$ in $U_{2}$, and we have for all $t \in(.50, .70]$

$$
c_{2} U_{2}=.20\left[\begin{array}{llll}
1 & 1 & 0 & 0 \\
0 & 0 & 0 & 1 \\
0 & 0 & 1 & 0
\end{array}\right]
$$

The sequence of transitions obtained in this fashion can be recorded in the matrix $\rho(U)$ with arrows and indices indicating the stage at which the transition occured. Thus we have at first in $\rho(U)$

$$
\rho(U)=\left[\begin{array}{cccc}
1 & 2 & 3 & 2 \\
2^{(2)} & \left.1^{( }\right) & 2 & 1 \\
3 & 3 & 1 & 3
\end{array}\right]
$$

Increasing $t$ above .70 results in degrading the membership of vector $\mathbf{x}_{\mathbf{1}}$ from class 1 to class 2 , as seen by first inspecting $\sigma(U)$ to locate the threshold, and then $\dot{\rho}(U)$ to find the path of the transition. This second switch is shown above in $\rho(U)$. Thus for all $t \in(.70, .80]$ we have

$$
c_{3} U_{3}=.10\left[\begin{array}{llll}
0 & 1 & 0 & 0 \\
1 & 0 & 0 & 1 \\
0 & 0 & 1 & 0
\end{array}\right]
$$

Sequential reclassification via the $R$-algorithm is perhaps most easily understood when $c=2$, for then each transition made merely moves one or more subjects from their maximum to minimum membership classes. To exemplify further, we rework the example of Section IV using the $R$-decomposition for matrix $U$ shown at (13):

$$
U=\left[\begin{array}{ccccc}
.90 & .80 & .30 & .40 & .05  \tag{13}\\
.10 & .20 & .70 & .60 & .95
\end{array}\right]
$$

From $U$ we find that

$$
\begin{aligned}
& \rho(U)=\left[\begin{array}{ccccc}
1 & 1 & 2 & 2 & 2 \\
2 & 2 & 1 & 1 & 1
\end{array}\right] ; \\
& \sigma(U)=\left[\begin{array}{ccccc}
.90 & .80 & 1 & 1 & 1 \\
1 & 1 & .70 & .60 & .95
\end{array}\right] ;
\end{aligned}
$$

$$
\begin{aligned}
& a_{0}=.00 \\
& a_{1}=.60>c_{1}=.60 \\
& a_{2}=.70>c_{2}=.10 \\
& a_{3}=.80>c_{3}=.10 \\
& a_{1}=.90>c_{4}=.10 \\
& a_{5}=.95>c_{5}=.05 \\
& a_{6}=1>c_{6}=.05
\end{aligned}
$$

To begin, we have from above that

$$
c_{1} U_{1}=.60\left[\begin{array}{lllll}
1 & 1 & 0 & 0 & 0 \\
0 & 0 & 1 & 1 & 1
\end{array}\right]
$$

with residual

$$
R_{1}=\left[\begin{array}{ccccc}
.30 & .20 & .30 & .40 & .05 \\
.10 & .20 & .10 & 0 & .35
\end{array}\right]
$$

While in a general decomposition the next coefficient $c_{2}=.10$ could be extracted in many ways, the $R$-algorithm proceeds by degrading the vector whose membership in its maximum membership class is first exceeded upon crossing the threshold $t=0.60$; in this case, from $\sigma(U)$ we find at address $(2,4)$ that $\mathbf{x}_{4}$ has the lowest tolerance, and must be reclassified (to address $(1,4)$ here) first:

$$
c_{2} U_{2}=.10\left[\begin{array}{lllll}
1 & 1 & 0 & 1 & 0 \\
0 & 0 & 1 & 0 & 1
\end{array}\right] .
$$

Increasing $t$ above .70 results in degrading the membership of $\mathbf{x}_{3}$

$$
c_{3} U_{3}=.10\left[\begin{array}{lllll}
1 & 1 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 1
\end{array}\right]
$$

At $t>.80, \sigma_{12}=.80 \Rightarrow$ reclassify $\mathbf{x}_{2}$ :

$$
c_{4} U_{4}=.10\left[\begin{array}{lllll}
1 & 0 & 1 & 1 & 0 \\
0 & 1 & 0 & 0 & 1
\end{array}\right]
$$

At $t<.90, \sigma_{11}=.90 \Rightarrow$ reclassify $\mathbf{x}_{1}$ :

$$
c_{5} U_{5}=.05\left[\begin{array}{lllll}
0 & 0 & 1 & 1 & 0 \\
1 & 1 & 0 & 0 & 1
\end{array}\right]
$$

Finally, for $t>.95, \sigma_{25}=.95$ requires us to reclassify $\mathbf{x}_{5}$ last; this is due to its most distinctive memberships in $U$ :

$$
c_{6} U_{6}=.05\left[\begin{array}{lllll}
0 & 0 & 1 & 1 & 1 \\
1 & 1 & 0 & 0 & 0
\end{array}\right]
$$

The entire sequence of transitions for this example may be recorded in the matrix $\rho(U)$ as follows:

From this we see that the sequence of transitions always proceeds "uphill" through the ranked columns of $\rho(U)$, from $1 \rightarrow 2 \cdots \rightarrow c$. Note that $U_{1}$ and the final partition $U_{6}$ are algebraic complements; it is always the case that the first and last partitions in an $R$-decomposition sequence are the maximum and minimum membership partitions which are the closest and farthest hard partitons associated with $U$ respectively. Since the $\left\{a_{k}\right\}$ were strictly ordered in this example, the $R$-decomposition was unique.

Since the first (and dominant) term of every $R$ and $M M$ decomposition is $U_{m m}$, one might conjecture that $U_{m m}$ is the dominant term in every convex decomposition of $U \in P_{f c}$. To see that this is not the case, consider the matrix

$$
U=\left[\begin{array}{llll}
.54 & .58 & .73 & .56  \tag{23}\\
.46 & .42 & .27 & .44
\end{array}\right]
$$

with maximum membership matrix $U_{m m}=\left[\begin{array}{ccc}1 & 1 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 0\end{array}\right]$. We exhibit for this $U$ the $M M$-decomposition, and a second convex decomposition which not only does not have $U_{m m}$ in it, but is in fact, entirely disjoint in the sense that no hard partition used in the $M M$-decomposition is used in the second one:
$M M$-decomposition of (23) Another disjoint decomposition

$$
\begin{array}{ll}
+.54\left[\begin{array}{llll}
1 & 1 & 1 & 1 \\
0 & 0 & 0 & 0
\end{array}\right] & +.44\left[\begin{array}{llll}
1 & 1 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right] \\
+.27\left[\begin{array}{llll}
0 & 0 & 0 & 0 \\
1 & 1 & 1 & 1
\end{array}\right] & +.29\left[\begin{array}{llll}
0 & 0 & 1 & 1 \\
1 & 1 & 0 & 0
\end{array}\right] \\
+.15\left[\begin{array}{llll}
0 & 0 & 1 & 0 \\
1 & 1 & 0 & 1
\end{array}\right] & +.13\left[\begin{array}{llll}
0 & 0 & 0 & 1 \\
1 & 1 & 1 & 0
\end{array}\right] \\
+.02\left[\begin{array}{llll}
0 & 1 & 1 & 1 \\
1 & 0 & 0 & 0
\end{array}\right] & +.10\left[\begin{array}{llll}
1 & 1 & 0 & 1 \\
0 & 0 & 1 & 0
\end{array}\right] \\
.02\left[\begin{array}{llll}
0 & 1 & 1 & 0 \\
1 & 0 & 0 & 1
\end{array}\right] . & +.04\left[\begin{array}{llll}
0 & 1 & 0 & 1 \\
1 & 0 & 1 & 0
\end{array}\right] .
\end{array}
$$

Although the term $U_{m m}$ which dominates both the $M M$ and $R$ decompositions does not always appear, it carries the coefficient $c_{1}$ which is maximum among all possible coefficients for a given $U$; accordingly, we presume that the equivalence
relation induced on $X \times X$ by $U_{m m}$ indicates the maximal strength of bonding at level $c_{1}$ which the vectors in $\Gamma$ enjoy when partitioned by $U$. A forthcoming paper will deal at length with the relationships between fuzzy partitions and fuzzy similarity relations in $X \times X$.

Finally, we observe that sequential application of the $R$ algorithm to its residuals at each step results in a $M M$-decomposition of $L^{\top}$.

## VI. Convex Decomposition and the Partition Coefficient

For $U \in V_{c n}$ the Euclidean matrix norm and and inner product from which it is derived are

$$
\begin{align*}
\|U\| & =\left(\operatorname{tr}\left(U U^{t}\right)\right)^{1 / 2}  \tag{24a}\\
\left\langle U, U^{\prime}\right\rangle & =\operatorname{tr}\left(U U^{t}\right)=\left\|L^{r}\right\|^{2} \tag{24b}
\end{align*}
$$

where $\operatorname{tr}(\cdot),(\cdot)^{t}$ respectively denote the trace and transpose operations. Following [5] we denote by $F: V_{c n} \rightarrow R$ the function

$$
\begin{equation*}
F(U)=\frac{\|U\|^{2}}{n} \tag{25}
\end{equation*}
$$

Since $F\left(U^{r}\right)$ is inversely proportional to the overall content of pairwise fuzzy intersections in $U$, it has been called the partition coefficient of $U$. It measures the overall "fuzziness" in $U$ via the following results proven in [5]:

$$
\begin{align*}
1 / c & \leqslant F(U)
\end{aligned} \leqslant 1 \forall U \in P_{f c}, ~ \begin{aligned}
1 / c & =F(\tilde{U}) \Leftrightarrow \tilde{U}=[1 / c]  \tag{26a}\\
1 & =F(U) \Leftrightarrow U \in P_{c 0} \text { is hard. } \tag{26b}
\end{align*}
$$

If we denote by $d(U, V)=\|U-V\|$ the distance between $U$ and $V$ in $V_{c n}$, it is easy to check that $d(U, \tilde{O})=(n(c-1) / c)^{1 / 2} \forall U \in P_{c 0}$; consequently, the matrix $\tilde{C}$ is the geometric centroid of $P_{f c}$. Loosely speaking, $\zeta$ is the fuzziest classification of the vectors in $X$ allowed in $P_{f c}$, whereas $U^{\prime} s \in P_{c 0}$ are the crispest (hardest).

Now suppose $U_{i}$ and $U_{j}$ are any pair of hard $c$-partitons of $X$. We shall write $U_{i} \leftrightarrow U_{3}$ to indicate that $U_{i}$ and $U_{j}$ are different classification states for X that can be realized from each other by a finite number $n_{i j}$ of transitions (or reclassifications) between stages $i$ and $j$. For example, if

$$
L_{\imath}=\left[\begin{array}{lllll}
1 & 1 & 0 & 0 & 0  \tag{27}\\
0 & 0 & 1 & 1 & 1
\end{array}\right] \quad \text { and } \quad U_{j}=\left[\begin{array}{ccccc}
1 & 0 & 0 & 1 & 0 \\
0 & 1 & 1 & 0 & 1
\end{array}\right]
$$

then $n_{i j}=2$ for $U_{i} \leftrightarrow U_{3}$. It is quite useful for our later work to note that $n_{i j}$ can be calculated readily as

$$
\begin{equation*}
n_{i j}=n-\operatorname{tr}\left(U_{i} U_{2}^{t}\right) \tag{28}
\end{equation*}
$$

To see this, consider number

$$
\begin{equation*}
N_{i j}=\operatorname{tr}\left(U_{i} U_{\jmath}^{t}\right)=\sum_{k=1}^{r} \sum_{t=1}^{n}\left(u_{i}\right)_{k t}\left(u_{j}\right)_{k t} \tag{29}
\end{equation*}
$$

Since $\left(u_{i}\right)_{k t}$ and $\left(u_{j}\right)_{k t}$ are respectively the membership of $\mathbf{x}_{t}$ in the $k$-th cluster of $U_{i}$ and $U_{j}$, each term in the sum at (29) is 1 iff $\mathbf{x}_{t}$ is not reclassified in the transition $U_{i} \leftrightarrow U_{j}$. Thus $N_{i j}$ is the total number of individuals not reclassified, and hence $n_{i j}=n-N_{i j}$ at (28) is the number which are reclassified, and

$$
\begin{equation*}
f_{2 j}=\frac{n_{\imath j}}{n}=\frac{n-N_{2 j}}{n} \tag{30}
\end{equation*}
$$

is the faction or percent of individuals reclassified during $U_{i} \leftrightarrow U_{j}$. For instance, we have for the matrices in (27) $n_{i j}=2, N_{\imath \jmath}=3$, and $f_{i j}=\frac{2}{5}=.40$, i.e., $40 \%$ of the data is reclassified.

The following theorem exhibits the relationship between the partition coefficient $F(U)$ and the amount of reclassification done in the terms of any convex decomposition of $U$ :

Theorem 4. Let $U \in P_{f c}, F(U)$ as in (25), $f_{i j}$ as in (29), and suppose that $U=\sum_{k=1}^{p} c_{k} U_{k}$ is any convex decomposition of $U$. Then

$$
\begin{equation*}
F(U)=1-\sum_{j=\jmath+1}^{p} \sum_{i=1}^{p-1} c_{i} f_{2 \nu} c_{j} . \tag{30}
\end{equation*}
$$

Proof.

$$
\begin{aligned}
F(U) & =F\left(\sum_{k=1}^{p} c_{k} U_{k}\right)=(1 / n)\left(\operatorname{tr}\left(\sum_{j=1}^{p} c_{j} U_{j}\right)\left(\sum_{j=1}^{p} c_{j} U_{j}\right)^{t}\right) \\
& =(1 / n)\left(\operatorname{tr}\left(\sum_{j=1}^{p} \sum_{i=1}^{p} c_{i} U_{i} U_{j} c_{c_{j}}\right)\right) \\
& =(1 / n)\left(\sum_{j=1}^{p} \sum_{i=1}^{p} c_{i} \operatorname{tr}\left(U_{i} U_{j}^{t}\right) c_{j}\right) \\
& =\sum_{j=1}^{p} \sum_{i=1}^{p} c_{i}\left(N_{\imath j} / n\right) c_{j}
\end{aligned}
$$

Observing that $N_{i j}=N_{j i} \forall i=\neq j$, and that $N_{u}=n F\left(U_{\imath}\right)=n \forall i$, we have

$$
\begin{aligned}
& F(U)=\sum_{j=1}^{p} c_{j}\left(N_{2 j} \mid n\right) c_{j}+2\left(\sum_{j=2+1}^{p} \sum_{l=1}^{p-1} c_{\imath}\left(N_{, j} \mid n\right) c_{j}\right) \\
& =\sum_{l=1}^{p} c_{l}^{2}+2\left(\sum_{j=\imath+1}^{p} \sum_{l=1}^{\nu-1} c_{\imath}\left(\left(n-\left(n-N_{\imath}\right) / n\right) c_{\jmath}\right)\right. \\
& =\left\{\sum_{i=1}^{p} c_{j}^{2}+2\left(\sum_{j=i+1}^{n} \sum_{i=1}^{x-1} c_{i} c_{j}\right)\right\}-2\left(\sum_{i=i+1}^{p} \sum_{i=1}^{n-1} c_{i} f_{i} c_{j}\right) \\
& =1-2\left(\sum_{j=2,1}^{p} \sum_{l=1}^{p-1} c_{i} f_{i j} c_{j}\right),
\end{aligned}
$$

the last equality following because the $c_{J}$ are convex coefficients, so they sum to 1 .
(2.E.D.

Equation (30) shows that $F(U)=1$ iff $f_{i j}=0 \forall i \neq j$ iff $n=-N_{i j} \forall i \neq j$ iff no point is ever reclassified in any convex decomposition of $U$ iff $U \in P_{c 0}$ is hard and is its own unique convex decomposition, $p=c_{1}=1$. An equivalent way to state Theorem 4 as as follows; we define the $p<p$ transition matrix

$$
\begin{equation*}
N=\left[\frac{N_{i j}}{n}\right] ; \quad N_{1 j}=\operatorname{tr}\left(U_{2} U_{3}^{t}\right) ; \quad 1 \therefore i, \quad j \therefore p \tag{31}
\end{equation*}
$$

Matrix $N$ is symmetric and positive-definite for $U \in P_{r 0}$, and may thus be used to induce the weighted inner product $\left\langle\mathbf{c}, \mathbf{c}{ }_{N}^{\prime}=\|\left.\mathbf{c}\right|_{N} ^{2}=\mathbf{c}^{t} N \mathbf{c}\right.$ on $\mathbb{R}^{\prime \prime}$. Lising this notation, we have from the proof of Theorem 4 the

Corollary. Let $U \in P_{f r}$, and suppose $U=\sum_{k=1}^{p} c_{k} U_{k}=-\sum_{k=1}^{\prime \prime} c_{k}^{\prime} U_{k}^{\prime}$ are taco convex decompositions of $U$. Let $N$ and $N^{\prime}$ be their respective transition matrices computed ria (31), and let $\mathbf{c}$ and $\mathbf{c}^{\prime}$ be their zectors of convex coefficients in $\mathbb{R}^{f}$ and $\mathbb{R}^{\text {q }}$. Then

$$
\begin{equation*}
\left.\mathbf{c}^{\prime}\right|_{N}=\left.\left.\right|^{\prime} \mathbf{c}^{\prime}\right|_{N^{\prime}} \tag{32}
\end{equation*}
$$

Proof. Expansion of $U$ in any decomposition leads to, for example, $F\left(l^{\prime}\right)$ $\sum_{1-1}^{p} \sum_{i-1}^{n}\left(c_{i} N_{i j} c_{j} \mid n\right)=\mathbf{c}^{t} N \mathbf{c}=|\mathbf{c}|_{N}^{2}$. Since $F\left(l^{V}\right)$ is fixed, (32) follows.
Q.E.D.

The implication of (32) is this; all convex decompositions of $C$ are equivalent in the sense that their vectors of convex coefficients have equal lengths when normalized by the amount of reclassification required by their respective sequences of hard transitions. Since the $i j$-th entry of $N$ is the fraction of individuals not reclassified in the transition $U_{2} \leftrightarrow U_{3}$, the value $F(U)=\mathbf{c}^{t} N \mathbf{c}$ maximizes (for different partitions $U$ ) when the overall amount of reclassification is
minimized: this is a new justification for the clustering strategy of maximizing $F$ outlined in [5]. Evidently the quadratic form $\mathbf{c}^{t} N \mathbf{c}$ increases as the number of "stable" points in the data does, maximizing when no reclassification is necessary.

Assuming without loss that $U_{i}$ and $U_{j}$ are distinct, it will be seen that $N_{i j} / n$ is always less than 1 for $i \neq j$. We emphasize here that this is the number or fraction of vectors not reclassified in the single-step transition $U_{i} \leftrightarrow U_{j}$, even though $U_{\imath}$ and $U_{3}$ may be actually realized via a multi-step transition which would accumulate a different fraction of fixed addresses. For example, the transition shown in (27) has $N_{i j} / n=3 / 5=.6$, so $60 \%$ of the data is stable during the transition. We could, however, realize $U_{3}$ from $U_{i}$ in the two-step transition

$$
U_{i}=\left[\begin{array}{lllll}
1 & 1 & 0 & 0 & 0  \tag{33}\\
0 & 0 & 1 & 1 & 1
\end{array}\right] \leftrightarrow U_{k}=\left[\begin{array}{lllll}
0 & 0 & 1 & 1 & 1 \\
1 & 1 & 0 & 0 & 0
\end{array}\right] \leftrightarrow U_{3}=\left[\begin{array}{lllll}
1 & 0 & 0 & 1 & 0 \\
0 & 1 & 1 & 0 & 1
\end{array}\right]
$$

In (33) we have $N_{i k} / n=0 / 5=.00 ; N_{k j} / n=2 / 5=.40$; and $N_{\imath j} / n=.60$ as above. Thus $U_{i} \leftrightarrow U_{j} \leftrightarrow U_{k i}$ but $N_{i j} \neq N_{\imath l i}+N_{k j}$.

To solidify these ideas we decompose the matrix at (34) with both the MM and $R$ algorithms:

$$
U=\left[\begin{array}{lllll}
.39 & .60 & .35 & .17 & .92  \tag{34}\\
.17 & .14 & .37 & .78 & .05 \\
.44 & .26 & .28 & .05 & .03
\end{array}\right] \doteq R_{0}
$$

(i) MM-decomposition of (34). The initial $M M$ path is $\mathbf{v}_{1}=\left(*_{31}, *_{12}, *_{23}\right.$, $*_{24}, *_{15}$ ), with minimum maximum membership $c_{1}=u_{23}=.37$,

$$
c_{1} U_{1}=.37\left[\begin{array}{ccccc}
0 & 1 & 0 & 0 & 1 \\
0 & 0 & 1 & 1 & 0 \\
1 & 0 & 0 & 0 & 0
\end{array}\right] ; \quad R_{1}=\left[\begin{array}{ccccc}
.39 & .23 & .35 & .17 & .55 \\
.17 & .14 & .00 & .41 & .05 \\
.07 & .26 & .28 & .05 & .03
\end{array}\right]
$$

From $\mathbf{v}_{2}=\left(*_{11}, *_{32}, *_{13}, *_{24}, *_{15}\right)$ we infer that $c_{2}=\left(r_{1}\right)_{32}=.26$, so

$$
\begin{array}{rlrl}
c_{2} U_{2}^{\top}=.26\left[\begin{array}{lllll}
1 & 0 & 1 & 0 & 1 \\
0 & 0 & 0 & 1 & 0 \\
0 & 1 & 0 & 0 & 0
\end{array}\right] ; & R_{2}=\left[\begin{array}{lllll}
.13 & .23 & .09 & .17 & .29 \\
.17 & .14 & .00 & .15 & .05 \\
.07 & .00 & .28 & .05 & .03
\end{array}\right] \\
\mathbf{v}_{\mathbf{3}}=\left(*_{21}, *_{12}, *_{33}, *_{14}, *_{15}\right) \Rightarrow c_{3}=\left(r_{2}\right)_{21}=\left(r_{2}\right)_{14}=.17, \text { from which } \\
c_{3} U_{3}=.17\left[\begin{array}{lllll}
0 & 1 & 0 & 1 & 1 \\
1 & 0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 & 0
\end{array}\right] ; & R_{3}=\left[\begin{array}{lllll}
.13 & .06 & .09 & .00 & .12 \\
.00 & .14 & .00 & .15 & .05 \\
.07 & .00 & .11 & .05 & .03
\end{array}\right] \\
\mathbf{v}_{4}=\left(*_{11}, *_{22}, *_{33}, *_{24}, *_{15}\right) \Rightarrow c_{4}=\left(r_{3}\right)_{33}=.11, \text { leading to } \\
c_{4} U_{4}=.11\left[\begin{array}{lllll}
1 & 0 & 0 & 0 & 1 \\
0 & 1 & 0 & 1 & 0 \\
0 & 0 & 1 & 0 & 0
\end{array}\right] ; & R_{4}=\left[\begin{array}{lllll}
.02 & .06 & .09 & .00 & .01 \\
.00 & .03 & .00 & .04 & .05 \\
.07 & .00 & .00 & .05 & .03
\end{array}\right]
\end{array}
$$

$\mathbf{v}_{5}=\left(*_{31}, *_{12}, *_{13}, *_{34}, *_{25}\right) \Rightarrow c_{5}=\left(r_{4}\right)_{34}=\left(r_{4}\right)_{25}=.05$, yielding

$$
c_{5} U_{5}=.05\left[\begin{array}{ccccc}
0 & 1 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 \\
1 & 0 & 0 & 1 & 0
\end{array}\right] ; \quad R_{5}=\left[\begin{array}{ccccc}
.02 & .01 & .04 & .00 & .01 \\
.00 & .03 & .00 & .04 & .00 \\
.02 & .00 & .00 & .00 & .03
\end{array}\right]
$$

At step 6 we encounter with $M M$-decomposition the first non-unique path of minimum maximum menbership in $R_{5}$ because of the tie in its first column. Choosing the upper path leads to $\mathbf{v}_{6}=\left(*_{11}, *_{22}, *_{13}, *_{24}, k_{35}\right)$ and the choice $c_{6}=\left(r_{5}\right)_{11}=.02$, whence

$$
c_{6} U_{6}=.02\left[\begin{array}{ccccc}
1 & 0 & 1 & 0 & 0 \\
0 & 1 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 1
\end{array}\right] ; \quad R_{6}=\left[\begin{array}{ccccc}
.00 & .01 & .02 & .00 & .01 \\
.00 & .01 & .00 & .02 & .00 \\
.02 & .00 & .00 & .00 & .01
\end{array}\right]
$$

In the seventh step there are four $M M$ paths in $R_{6}$. For the choice $\mathbf{v}_{7}=\left(*_{31}, *_{12}, *_{13}, *_{24}, *_{15}\right), c_{7}=\left(r_{6}\right)_{12}=\left(r_{6}\right)_{15}$ gives

$$
c_{7} U_{7}-.01\left[\begin{array}{ccccc}
0 & 1 & 1 & 0 & 1 \\
0 & 0 & 0 & 1 & 0 \\
1 & 0 & 0 & 0 & 0
\end{array}\right] ; \quad R_{7}=\left[\begin{array}{ccccc}
.00 & .00 & .01 & .00 & .00 \\
.00 & .01 & .00 & .01 & .00 \\
.01 & .00 & .00 & .00 & .01
\end{array}\right]
$$

and from $R_{7}$ we have at the last stage of $M M$-decomposition

$$
c_{8} U_{8}=.01\left[\begin{array}{lllll}
0 & 0 & 1 & 0 & 0 \\
0 & 1 & 0 & 1 & 0 \\
1 & 0 & 0 & 0 & 1
\end{array}\right] ;
$$

with $R_{8}=\theta$, the zero matrix.
(ii) $R$-Decomposition of (34). This decomposition begins with the construction of the ranking and cumulative sum matrices $\rho(U)$ and $\sigma(U)$. Repeating (34) for convenience, we have

$$
\begin{aligned}
U & =\left[\begin{array}{ccccc}
.39 & .60 & .35 & .17 & .92 \\
.17 & .14 & .37 & .78 & .05 \\
.44 & .26 & .28 & .05 & .03
\end{array}\right] ; \\
\rho(U) & =\left[\begin{array}{ccccc}
2 & 1 & 2 & 2 & 1 \\
3 & 3 & 1 & 1 & 2 \\
1 & 2 & 3 & 3 & 3
\end{array}\right] ; \\
\sigma(U) & =\left[\begin{array}{ccccc}
.83 & .60 & .72 & .95 & .92 \\
1 & 1 & .37 & .78 & .97 \\
.44 & .86 & 1 & 1 & 1
\end{array}\right] .
\end{aligned}
$$

From $\sigma(U)$ we array the thresholds $\left\{a_{k}\right\}$ and from them derive the successive differences, say $\left\{d_{k}\right\}$ :

$$
\begin{aligned}
& a_{0}=.00>d_{1}=.37 \\
& a_{1}=.37>d_{2}=.07 \\
& a_{2}-.44>d_{3}=.16 \\
& a_{3}=.60>d_{4}=.12 \\
& a_{4}=.72>d_{5}=.06 \\
& a_{5}=.78>d_{6}=.05 \\
& a_{6}=.83 \\
& a_{7}=.86>d_{7}=.03 \\
& a_{8}=.92>d_{8}=.06 \\
& a_{9}=.95>d_{9}=.03 \\
& a_{10}=.97>d_{10}=.02 \\
& a_{11}=1>d_{11}=.03 .
\end{aligned}
$$

Using $\sum d_{k} W_{k}$ for the form of the $R$-decomposition, the initial term $d_{1} W_{1}$ is exactly the same as $c_{1} U_{1}$ in the $M M$-decomposition above, with identical residual. Upgrading the threshold from .00 , all data points are stable until $t>.37$, at which time $\mathbf{x}_{3}$ must be transferred from class 2 to class 1 (from $\sigma(U)$ we see that $\mathbf{x}_{3}$ will now remain in class 1 until $t>.72$, at which time it will transfer from class 1 to class 3 ). Reclassifying $\mathbf{x}_{3}$ results in the second term

$$
d_{2} W_{2}=.07\left[\begin{array}{lllll}
0 & 1 & 1 & 0 & 1 \\
0 & 0 & 0 & 1 & 0 \\
1 & 0 & 0 & 0 & 0
\end{array}\right] ; \quad R_{2}=\left[\begin{array}{ccccc}
.39 & .16 & .28 & .17 & .48 \\
.17 & .14 & .00 & .34 & .05 \\
.00 & .26 & .28 & .05 & .03
\end{array}\right] .
$$

The threshold is now increased without change in classification until it crosses the next cumulative sum, $t>.44$. At this stage $\mathrm{x}_{1}$ becomes unstable in class 3 and must be transferred to class 1 , where it will remain until $t>.83$;

$$
d_{3} W_{3}=.16\left[\begin{array}{ccccc}
1 & 1 & 1 & 0 & 1 \\
0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0
\end{array}\right] ; \quad R_{3}=\left[\begin{array}{ccccc}
.23 & .00 & .12 & .17 & .32 \\
.17 & .14 & .00 & .18 & .05 \\
.00 & .26 & .28 & .05 & .03
\end{array}\right]
$$

Since the ordering of the $\left\{a_{k}\right\}$ was strict, we expect a unique 11 term $R$-decomposition of $U$. To shorten this example we record the unique sequence of reclassifications as address transfers in the matrix $\rho(U)$ :


The sequence of transitions recorded in $\rho(U)$ indicates the stability of each of the $W_{k}$ 's to cumulative membership thresholding. The sequence suggests that $\mathbf{x}_{3}$ is the least stable or most undistinctive data point, since it is transferred twice in the early stages of decomposition, whereas $\mathbf{x}_{5}$ holds its initial classification the longest, indicating its relative certainty of membership distribution in the given fuzzy partition.

The partition coefficient $F(U)$ for this $U$ is $F(U)=\mathbf{c}^{t} N \mathbf{c}=\mathbf{d}^{t} N^{\prime} \mathbf{d}=$ $\operatorname{tr}\left(U U^{t}\right) / n=.53$, where

$$
\begin{aligned}
& \mathbf{c}^{t}=(.37, .26, .17, .11, .05, .02, .01, .01) \\
& \mathbf{d}^{t}=(.37, .07, .16, .12, .06, .05, .03, .06, .03, .02, .03)
\end{aligned}
$$

and $N, N^{\prime}$ are the $8 \times 8$ and $11 \times 11$ transition matrices calculated via (31) for the $M M$ and $R$ decompositions respectively. Since $c=3$ we know that $.33 \leqslant F(U) \leqslant 1$ : accordingly, .53 is a relatively low partition coefficient, indicating a fuzzy partition containing quite a bit of fuzzy uncertainty. This is manifested in the convex decompositions by coefficients which have no strongly dominant term. Thus $c_{1} U_{1}=d_{1} W_{1}=.37 U_{m m}$ indicates a relatively low minimum maximum membership bonding for the relation induced on the data by $U_{m m}$; indeed, for $c=3$ the weakest possible dominant term has for its coefficient .33, so that one's confidence (of .37) in taking $U_{m m}{ }_{m}$ as a first approximation to hard clusters in $X$ should be slight.

## VII. Summary

Several important facts about convex decompositions of fuzzy c-partitions of X have been enumerated. In particular, the dimension of fuzzy $c$-partition space is $n(c-1)$. Two algorithms for decomposition have been analyzed and exemplified. Both of these algorithms are rather tedious to work with by hand, but are quite readily programmed for large scale problems involving real data sets. The relationship between convex decomposition and the partition coefficient provides additional justification for a previously advocated fuzzy clustering strategy. The authors plan to devote a forthcoming paper to the connection between fuzzy partition spaces and fuzzy similarity relation spaces, in which analogs of the convex decompositions descrihed above will hopefully lead to a new method for clustering with fuzzy graphs.

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